

# Solving Dynamic Games with Newton's Method

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July 31, 2008

## Discrete-Time Finite-State Stochastic Games

Central tool in analysis of strategic interactions among forward-looking players in dynamic environments

Example: The Ericson & Pakes (1995) model of dynamic competition in an oligopolistic industry

Little analytical tractability

Most popular tool in the analysis: The Pakes & McGuire (1994) algorithm to solve numerically for an MPE (and variants thereof)

## Applications

Advertising (Doraszelski & Markovich 2007)

Capacity accumulation (Besanko & Doraszelski 2004, Chen 2005, Ryan 2005, Beresteanu & Ellickson 2005)

Collusion (Fershtman & Pakes 2000, 2005, de Roos 2004)

Consumer learning (Ching 2002)

Firm size distribution (Laincz & Rodrigues 2004)

Learning by doing (Benkard 2004, Besanko, Doraszelski, Kryukov & Satterthwaite 2007)

## Applications cont'd

Mergers (Berry & Pakes 1993, Gowrisankaran 1999)

Network externalities (Jenkins, Liu, Matzkin & McFadden 2004, Markovich 2004, Markovich & Moenius 2007)

Productivity growth (Laincz 2005)

R&D (Gowrisankaran & Town 1997, Auerswald 2001, Song 2002, Yeltekin et al. 2007)

Technology adoption (Schivardi & Schneider 2005)

International trade (Erdem & Tybout 2003)

Finance (Goettler, Parlour & Rajan 2004, Kadyrzhanova 2005).

## Need for better Computational Techniques

Doraszelski and Pakes (2007, in: Handbook of IO)

“Moreover the burden of currently available techniques for computing equilibria to the models we do know how to analyze is still large enough to be a limiting factor in the analysis of many empirical and theoretical issues of interest.”

## This Tutorial

1. Discrete-Time Finite-State Stochastic Games
2. Separable Game
3. Solution Methods for Dynamic Games

# Discrete-Time Finite-Space Stochastic Games

## State Space

Infinite-horizon game in discrete time  $t = 0, 1, 2, \dots$

Set of  $N$  players,  $i = 1, \dots, N$

At time  $t$  player  $i$  is in one of finitely many states  $x_t^i \in X^i$

State space of the game  $X = \prod_i X^i$

State in period  $t$  is  $x_t = (x_t^1, \dots, x_t^N)$

Notation:  $x_t^{-i} = (x_t^1, \dots, x_t^{i-1}, x_t^{i+1}, \dots, x_t^N)$

## Player's Actions and Transitions

Player  $i$ 's action in period  $t$  is  $u_t^i \in U^i(x_t)$

Set of feasible actions  $U^i(x_t)$  is arbitrary, often  $U^i = \mathbb{R}_+^K$

Players' actions at time  $t$ :  $u_t = (u_t^1, \dots, u_t^N)$

Law of motion: State follows a controlled discrete-time, finite-state, first-order Markov process with transition probability  $\Pr(x'|u_t, x_t)$

Special case of independent transitions:

$$\Pr(x'|u_t, x_t) = \prod_{i=1}^N \Pr^i\left((x')^i | u_t^i, x_t^i\right)$$

## Objective Function

Player  $i$  receives a payoff of  $\pi^i(u_t, x_t)$  in period  $t$

Objective is to maximize the expected NPV of future cash flows

$$E \left\{ \sum_{t=0}^{\infty} \beta^t \pi^i(u_t, x_t) \right\},$$

with discount factor  $\beta \in (0, 1)$

## Bellman Equation

$V^i(x)$  is the expected NPV to player  $i$  if the current state is  $x$

Bellman equation for player  $i$  is

$$V^i(x) = \max_{u^i} \pi^i(u^i, U^{-i}(x), x) + \beta \mathbb{E}_{x'} \{V^i(x') | u^i, U^{-i}(x), x\} \quad (1)$$

where  $U^{-i}(x)$  denotes feedback (Markovian) strategies of other players

Player  $i$ 's strategy is given by

$$U^i(x) = \arg \max_{u^i} \pi^i(u^i, U^{-i}(x), x) + \beta \mathbb{E}_{x'} \{V^i(x') | u^i, U^{-i}(x), x\} \quad (2)$$

System of equations defined by (1) and (2) for each player  $i = 1, \dots, N$  and each state  $x \in X$  defines a pure-strategy MPE

## Example of a Separable Game: Patent Race

## Patent Race Between Two Firms

$N$  innovation stages

Firms start race at stage 0

Period  $t$  innovation stages:  $(x_{1,t}, x_{2,t})$  where  
 $x_{i,t} \in X \equiv \{0, \dots, N\}, i = 1, 2$

Period  $t$  investment:  $a_{i,t} \in A = [0, \bar{A}] \subset \mathbb{R}_+, i = 1, 2$

Cost of investment:  $C_i(a) = c_i a^\eta, \eta \in \mathbb{N}, c_i > 0, i = 1, 2$

Independent and stochastic innovation technologies

## Transition from State to State

Transition from period to period:  $x_{i,t+1} = x_{i,t}$  or  $x_{i,t+1} = x_{i,t} + 1$

Markov process (depends on investment levels)

Firm  $i$ 's state evolves according to

$$x_{i,t+1} = \begin{cases} x_{i,t}, & \text{with probability } p(x_{i,t}|a_{i,t}, x_{i,t}) \\ x_{i,t} + 1, & \text{with probability } p(x_{i,t} + 1|a_{i,t}, x_{i,t}) \end{cases}$$

Distribution over next period's states

$$p(x|a, x) = F(x|x)^a$$

$$p(x + 1|a, x) = 1 - F(x|x)^a$$

$F(x|x) \in (0, 1)$  is probability that there is no change in state if  $a = 1$

## Firms' Optimization Problem

First firm to reach state  $N$  wins the race and receives prize  $\Omega$

Ties are broken by flip of a coin

Firms discount future costs and revenues at common rate  $\beta < 1$

Firms' objective: maximize expected discounted payoffs

## Equilibrium I

Restriction to pure Markov strategies

Firm  $i$ 's strategy:  $\sigma_i(\cdot) : X \times X \rightarrow A$

Expected discounted payoff:  $V_i(\cdot) : X \times X \rightarrow \mathbb{R}$

Bellmann equation for  $x_i, x_{-i} < N$ ,

$$V_i(x_i, x_{-i}) =$$

$$\max_{a_i \in A} \left\{ -C_i(a_i) + \beta \sum_{x'_i, x'_{-i}} p(x'_i | a_i, x_i) p(x'_{-i} | a_{-i}, x_{-i}) V_i(x'_i, x'_{-i}) \right\}$$

## Equilibrium II

Boundary condition at terminal states

$$V_i(x_i, x_{-i}) = \begin{cases} \Omega, & \text{for } x_{-i} < x_i = N \\ \Omega/2, & \text{for } x_i = x_{-i} = N \\ 0, & \text{for } x_i < x_{-i} = N \end{cases}$$

Optimal strategies satisfy

$$\sigma_i(x_i, x_{-i}) =$$

$$\arg \max_{a_i \in A} \left\{ -C_i(a_i) + \beta \sum_{x'_i, x'_{-i}} p(x'_i | a_i, x_i) p(x'_{-i} | a_{-i}, x_{-i}) V_i(x'_i, x'_{-i}) \right\}$$

## Our Equilibrium Equations

$$0 = -V_i(x_i, x_{-i}) - c_i a_i^\eta + \beta \sum_{x'_i, x'_{-i}} p(x'_i | a_i, x_i) p(x'_{-i} | a_{-i}, x_{-i}) V_i(x'_i, x'_{-i})$$

$$0 = -\eta c_i a_i^{\eta-1} + \beta \sum_{x'_i, x'_{-i}} \frac{\partial}{\partial a_i} p(x'_i | a_i, x_i) p(x'_{-i} | a_{-i}, x_{-i}) V_i(x'_i, x'_{-i})$$

Parameter specification:  $c_1, c_2, \eta, F(x_1, x_2) \equiv F, \Omega$

Unknowns:  $V_1(x_1, x_2), V_2(x_1, x_2), a_1(x_1, x_2), a_2(x_1, x_2)$

Four equations per stage  $(x_i, x_{-i})$

Backward induction: instead of solving all equations simultaneously

solve each stage game separately

# Solving Systems of Nonlinear Equations

## Nonlinear Systems of Equations

System  $F(x) = 0$  of  $n$  nonlinear equations in  $n$  variables  
 $x = (x_1, x_2, \dots, x_n) \in \mathbb{R}^n$

$$\begin{aligned} F_1(x_1, x_2, \dots, x_n) &= 0 \\ F_2(x_1, x_2, \dots, x_n) &= 0 \\ &\vdots \\ F_{n-1}(x_1, x_2, \dots, x_n) &= 0 \\ F_n(x_1, x_2, \dots, x_n) &= 0 \end{aligned}$$

Initial guess  $x^0 = (x_1^0, x_2^0, \dots, x_n^0)$

Methods generate a sequence of iterates  $x^0, x^1, x^2, \dots, x^k, x^{k+1}, \dots$

## Solution Methods

Most popular methods in economics for solving  $F(x) = 0$

1. Gauss-Jacobi Method
2. Gauss-Seidel Method
3. Newton's Method
4. Homotopy Methods

## Gauss-Jacobi Method

Last iterate  $x^k = (x_1^k, x_2^k, x_3^k, \dots, x_{n-1}^k, x_n^k)$

New iterate  $x^{k+1}$  computed by repeatedly solving one equation in one variable using only values from  $x^k$

$$F_1(x_1^{k+1}, x_2^k, x_3^k, \dots, x_{n-1}^k, x_n^k) = 0$$

$$F_2(x_1^k, x_2^{k+1}, x_3^k, \dots, x_{n-1}^k, x_n^k) = 0$$

$$\vdots$$

$$F_{n-1}(x_1^k, x_2^k, \dots, x_{n-2}^k, x_{n-1}^{k+1}, x_n^k) = 0$$

$$F_n(x_1^k, x_2^k, \dots, x_{n-2}^k, x_{n-1}^k, x_n^{k+1}) = 0$$

Computer storage: Need to store both  $x^k$  and  $x^{k+1}$

Interpretation as iterated simultaneous best reply

## Gauss-Seidel Method

Last iterate  $x^k = (x_1^k, x_2^k, x_3^k, \dots, x_{n-1}^k, x_n^k)$

New iterate  $x^{k+1}$  computed by repeatedly solving one equation in one variable and immediately updating the iterate

$$\begin{aligned} F_1(x_1^{k+1}, x_2^k, x_3^k, \dots, x_{n-1}^k, x_n^k) &= 0 \\ F_2(x_1^{k+1}, x_2^{k+1}, x_3^k, \dots, x_{n-1}^k, x_n^k) &= 0 \\ &\vdots \\ F_{n-1}(x_1^{k+1}, x_2^{k+1}, \dots, x_{n-2}^{k+1}, x_{n-1}^{k+1}, x_n^k) &= 0 \\ F_n(x_1^{k+1}, x_2^{k+1}, \dots, x_{n-2}^{k+1}, x_{n-1}^{k+1}, x_n^{k+1}) &= 0 \end{aligned}$$

Computer storage: Need to store only one vector

Interpretation as iterated sequential best reply

## Solving a Simple Cournot Game

$N$  firms

Firm  $i$ 's production quantity  $q_i$

Total output is  $Q = q_1 + q_2 + \dots + q_N$

Linear inverse demand function,  $P(Q) = A - Q$

All firms have identical cost functions  $C(q) = \frac{2}{3}cq^{3/2}$

Firm  $i$ 's profit function  $\Pi_i$  is

$$\Pi_i = q_i P(q_i + Q_{-i}) - C(q_i) = q_i (A - (q_i + Q_{-i})) - \frac{2}{3}cq_i^{3/2}$$

## First-order Conditions

Necessary and sufficient first-order conditions

$$A - Q_{-i} - 2q_i - c\sqrt{q_i} = 0$$

Firm  $i$ 's best reply  $R(Q_{-i})$  to a production quantity  $Q_{-i}$  of its competitors

$$q_i = R(Q_{-i}) = \left( \frac{A - Q_{-i}}{2} + \frac{c^2}{8} \right) - \frac{c}{2} \sqrt{\frac{A - Q_{-i}}{2} + \frac{c^2}{16}}$$

Parameter values:  $N = 2$  firms,  $A = 145$ ,  $c = 4$

## Solving the Cournot Game with Gauss-Jacobi

$k$	$q_i^k$	$\max_i  q_i^k - q_i^{k-1} $
0	10	—
1	52.9471	42.9471
2	34.3113	18.6358
3	42.3318	8.02047
4	38.8656	3.46613
5	40.3611	1.49545
6	39.7154	0.645682
7	39.9941	0.278695
15	39.9102	0.000336014
16	39.9100	0.000145047
20	39.910075	5.036 (-6)
21	39.910078	2.174 (-6)

## Solving the Cournot Game with Gauss-Seidel

$k$	$q_1^k$	$q_2^k$	$\max_i  q_i^k - q_i^{k-1} $
0	10	10	—
1	52.9471	34.3113	42.9471
2	42.3318	38.8656	10.6153
3	40.3611	39.7154	1.97068
4	39.9941	39.8738	0.366987
5	39.9257	39.9033	0.0683762
6	39.913	39.9088	0.0127409
7	39.9106	39.9098	0.00237412
8	39.9102	39.91	0.000442391
9	39.9101	39.9101	0.0000824347
10	39.9101	39.9101	0.0000153608
11	39.91008	39.91008	2.862 (-6)

## Gauss-Jacobi with $N = 4$ firms blows up

Cournot equilibrium  $q^i = 25$  for all firms

$$x^0 = (24, 25, 25, 25)$$

$k$	$q_1^k$	$q_2^k = q_3^k = q_4^k$	$\max_i  q_i^k - q_i^{k-1} $
1	25	25.4170	1
2	24.4793	24.6527	0.7642
3	25.4344	25.5068	0.9551
4	24.3672	24.3973	1.1095
5	25.7543	25.7669	1.3871
13	29.5606	29.5606	8.1836
14	19.3593	19.3593	10.201
15	32.1252	32.1252	12.766
20	4.8197	4.8197	37.373
21	50.9891	50.9891	46.169

## Newton's Method

Foundation of Newton's Method: **Taylor's Theorem**

**THEOREM.** Suppose the function  $F : X \rightarrow \mathbb{R}^m$  is continuously differentiable on the open set  $X \subset \mathbb{R}^n$  and that the Jacobian function  $J_F$  is Lipschitz continuous at  $x$  with Lipschitz constant  $\gamma^l(x)$ . Also suppose that for  $s \in \mathbb{R}^n$  the line segment  $x + \theta s \in X$  for all  $\theta \in [0, 1]$ . Then, the linear function  $L(s) = F(x) + J_F(x)s$  satisfies

$$\|F(x + s) - L(s)\| \leq \frac{1}{2} \gamma^l(x) \|s\|^2 .$$

Taylor's Theorem suggests the approximation

$$F(x + s) \approx L(s) = F(x) + J_F(x)s$$

## Newton's Method in Pure Form

Initial guess  $x^0$

Given iterate  $x^k$  choose Newton step by calculating a solution  $s^k$  to the system of linear equations

$$J_F(x^k) s^k = -F(x^k)$$

New iterate  $x^{k+1} = x^k + s^k$

Excellent local convergence properties

Solving Cournot Game ( $N = 4$ ) with Newton's Method

$k$	$q_i^k$	$\max_i  q_i^k - q_i^{k-1} $
0	10	—
1	24.6208	14.6208
2	24.9999	0.3791
3	25.0000	0.000108
4	25.0000	8.67(-12)

## Shortcomings of Newton's Method

If initial guess  $x^0$  is far from a solution Newton's method may behave erratically; for example, it may diverge or cycle (!)

If  $J_F(x^k)$  is singular the Newton step may not be defined

It may be too expensive to compute the Newton step  $s^k$  for large systems of equations

The root  $x^*$  may be degenerate ( $J_F(x^*)$  is singular) and convergence is very slow

Practical variants of Newton-like methods overcome all these issues

## Practical Newton-like Method

General idea: Obtain global (!) convergence by combining the Newton step with line-search or trust-region methods from optimization

Merit function monitors progress towards root of  $F$

Most widely used merit function is sum of squares

$$M(x) = \frac{1}{2} \|F(x)\|^2 = \frac{1}{2} \sum_{i=1}^n F_i^2(x)$$

Any root  $x^*$  of  $F$  yields global minimum of  $M$

Local minimizers with  $M(x) > 0$  are not roots of  $F$

$$\nabla M(\tilde{x}) = J_F(\tilde{x})^\top F(\tilde{x}) = 0$$

and so  $F(\tilde{x}) \neq 0$  implies  $J_F(\tilde{x})$  is singular

## Line Search Method

Newton step

$$J_f(x^k) s^k = -F(x^k)$$

yields a descent direction of  $M$  as long as  $F(x^k) \neq 0$

$$(s^k)^\top \nabla M(x^k) = (s^k)^\top J_F(x^k)^\top F(x^k) = -\|F(x^k)\|^2 < 0$$

Given step length  $\alpha^k$  the new iterate is

$$x^{k+1} = x^k + \alpha^k s^k$$

## Step length

Inexact line search condition (Armijo condition)

$$M(x^k + \alpha s^k) \leq M(x^k) + c \alpha \left( \nabla M(x^k) \right)^T s^k$$

for some constant  $c \in (0, 1)$

Step length is the largest  $\alpha$  satisfying the inequality

For example, try  $\alpha = 1, \frac{1}{2}, \frac{1}{2^2}, \frac{1}{2^3}, \dots$

This approach is not Newton's method for minimization

No computation or storage of Hessian matrix

## Global Convergence Property

**THEOREM.** Suppose that  $J_F$  is Lipschitz continuous and both  $\|J_F(x)\|$  and  $\|F(x)\|$  are bounded above in an open neighborhood of the level set  $\{x : M(x) \leq M(x^0)\}$ . Under some further mild technical conditions the sequence of iterates  $x^0, x^1, \dots, x^k, x^{k+1}, \dots$  satisfies

$$\left(J_F(x^k)\right)^\top F(x^k) \rightarrow 0$$

as  $k \rightarrow \infty$ .

Moreover, if  $\|J_F(x^k)\| \geq \delta > 0$  then

$$F(x^k) \rightarrow 0.$$

## Cournot Game with Learning and Investment

$N = 2$  firms in dynamic Cournot competition

State of the game: production cost of two firms

Each period: Firms engage in quantity competition

Stochastic transition to state in next period depends on three forces

Learning: Current output may lead to lower production cost

Investment: Firms can also make investment expenditures to reduce cost

Depreciation: Shock to efficiency may increase cost

## Period Game

Firm  $i$ 's production quantity  $q_i$

Total output is  $Q = q_1 + q_2$

Linear inverse demand function,  $P(Q) = A - Q$

Firms' production cost functions are quadratic  $CP_i(q) = \frac{1}{2}b_iq^2$

Firms' profit functions are

$$\Pi_1 = q_1 P(q_1 + q_2) - \theta_1 \left( \frac{1}{2}b_1q_1^2 \right)$$

$$\Pi_2 = q_2 P(q_1 + q_2) - \theta_2 \left( \frac{1}{2}b_2q_2^2 \right)$$

Efficiency of firm  $i$  is given by  $\theta_i$

## Dynamic Setting

Each firm can be in one of  $S$  states,  $j = 1, 2, \dots, S$

State  $j$  of firm  $i$  determines its efficiency level

$\theta_i = \Theta^{(j-1)/(S-1)}$  for some  $\Theta \in (0, 1)$

Total range of efficiency levels  $[\Theta, 1]$  for any  $S$

Possible transitions from state  $j$  to states  $j - 1, j, j + 1$  in next period

Transition probabilities for firm  $i$  depend on  
production quantity  $q_i$

investment effort  $u_i$

depreciation shock

## Transition Probabilities

Probability of successful learning ( $j$  to  $j + 1$ ),  $\psi(q) = \frac{\kappa q}{1 + \kappa q}$

Probability of successful investment ( $j$  to  $j + 1$ ),  $\phi(u) = \frac{\alpha u}{1 + \alpha u}$

Cost of investment for firm  $i$ ,  $CI_i(u) = \frac{1}{S-1} \left( \frac{1}{2} d_i u^2 \right)$

Probability of depreciation shock,  $\delta$

These individual probabilities, appropriately combined, yield transition probabilities

## Equilibrium Equations

Bellman equation for each firm

First-order condition w.r.t. quantity  $q_i$

First-order condition w.r.t. investment  $u_i$

Three equations per firm per state

Total of 6  $S^2$  equations

## GAMS Code I

$$\begin{aligned}
 V1(m1e,m2e) = & e = Q1(m1e,m2e)*(1 - Q1(m1e,m2e)/M - \\
 & Q2(m1e,m2e)/M) - ((b1*power(Q1(m1e,m2e),2))/2. + \\
 & a1*Q1(m1e,m2e))*theta1(m1e) - ((d1*power(U1(m1e,m2e),2))/2. + \\
 & c1*U1(m1e,m2e))/(-1 + Nst) + (beta*((1 - 2*delta + power(delta,2) \\
 & + Q2(m1e,m2e)*(delta*kappa - kappa*power(delta,2) + \\
 & alpha*kappa*power(delta,2)*U1(m1e,m2e)) + (alpha*delta - \\
 & alpha*power(delta,2))*U2(m1e,m2e) + Q1(m1e,m2e)*(delta*kappa - \\
 & kappa*power(delta,2) + power(delta,2)*power(kappa,2)*Q2(m1e,m2e) \\
 & + alpha*kappa*power(delta,2)*U2(m1e,m2e)) + \\
 & U1(m1e,m2e)*(alpha*delta - alpha*power(delta,2) +
 \end{aligned}$$

## GAMS Code II

```

power(alpha,2)*power(delta,2)*U2(m1e,m2e))) * V1(m1e,m2e) + (delta -
power(delta,2) + kappa*power(delta,2)*Q1(m1e,m2e) +
alpha*power(delta,2)*U1(m1e,m2e)) * V1(m1e,m2e - 1) + ((alpha -
2*alpha*delta + alpha*power(delta,2)) * U2(m1e,m2e) +
(delta*power(alpha,2) -
power(alpha,2)*power(delta,2)) * U1(m1e,m2e) * U2(m1e,m2e) +
Q2(m1e,m2e) * (kappa - 2*delta*kappa + kappa*power(delta,2) +
(alpha*kappa - alpha*delta*kappa) * U2(m1e,m2e) +
U1(m1e,m2e) * (alpha*delta*kappa - alpha*kappa*power(delta,2) +
delta*kappa*power(alpha,2) * U2(m1e,m2e))) +
Q1(m1e,m2e) * ((alpha*delta*kappa -

```

## GAMS Code III

$$\begin{aligned}
 & \alpha * \kappa * \text{power}(\text{delta}, 2) * U2(m1e, m2e) + \\
 & Q2(m1e, m2e) * (\text{delta} * \text{power}(\kappa, 2) - \text{power}(\text{delta}, 2) * \text{power}(\kappa, 2) \\
 & + \alpha * \text{delta} * \text{power}(\kappa, 2) * U2(m1e, m2e))) * V1(m1e, m2e + 1) + \\
 & (\text{delta} - \text{power}(\text{delta}, 2) + \kappa * \text{power}(\text{delta}, 2) * Q2(m1e, m2e) + \\
 & \alpha * \text{power}(\text{delta}, 2) * U2(m1e, m2e)) * V1(m1e - 1, m2e) + \\
 & \text{power}(\text{delta}, 2) * V1(m1e - 1, m2e - 1) + ((\alpha * \text{delta} - \\
 & \alpha * \text{power}(\text{delta}, 2)) * U2(m1e, m2e) + Q2(m1e, m2e) * (\text{delta} * \kappa - \\
 & \kappa * \text{power}(\text{delta}, 2) + \alpha * \text{delta} * \kappa * U2(m1e, m2e))) * V1(m1e - \\
 & 1, m2e + 1) + ((\alpha * \text{delta} * \kappa - \\
 & \alpha * \kappa * \text{power}(\text{delta}, 2)) * Q2(m1e, m2e) * U1(m1e, m2e) + \\
 & U1(m1e, m2e) * (\alpha - 2 * \alpha * \text{delta} + \alpha * \text{power}(\text{delta}, 2) + \\
 & (\text{delta} * \text{power}(\alpha, 2) -
 \end{aligned}$$

## GAMS Code IV

```

power(alpha,2)*power(delta,2)*U2(m1e,m2e)) + Q1(m1e,m2e)*(kappa
- 2*delta*kappa + kappa*power(delta,2) +
Q2(m1e,m2e)*(delta*power(kappa,2) - power(delta,2)*power(kappa,2)
+ alpha*delta*power(kappa,2)*U1(m1e,m2e)) + (alpha*delta*kappa -
alpha*kappa*power(delta,2))*U2(m1e,m2e) +
U1(m1e,m2e)*(alpha*kappa - alpha*delta*kappa +
delta*kappa*power(alpha,2)*U2(m1e,m2e))))*V1(m1e + 1,m2e) +
((alpha*delta - alpha*power(delta,2))*U1(m1e,m2e) +
Q1(m1e,m2e)*(delta*kappa - kappa*power(delta,2) +
alpha*delta*kappa*U1(m1e,m2e)))*V1(m1e + 1,m2e - 1) +
((power(alpha,2) - 2*delta*power(alpha,2) +
power(alpha,2)*power(delta,2))*U1(m1e,m2e)*U2(m1e,m2e) +

```

## GAMS Code V

$$\begin{aligned}
 & Q2(m1e,m2e)*U1(m1e,m2e)*(alpha*kappa - 2*alpha*delta*kappa + \\
 & alpha*kappa*power(delta,2) + (kappa*power(alpha,2) - \\
 & delta*kappa*power(alpha,2))*U2(m1e,m2e)) + \\
 & Q1(m1e,m2e)*((alpha*kappa - 2*alpha*delta*kappa + \\
 & alpha*kappa*power(delta,2))*U2(m1e,m2e) + (kappa*power(alpha,2) - \\
 & delta*kappa*power(alpha,2))*U1(m1e,m2e)*U2(m1e,m2e) + \\
 & Q2(m1e,m2e)*(power(kappa,2) - 2*delta*power(kappa,2) + \\
 & power(delta,2)*power(kappa,2) + (alpha*power(kappa,2) - \\
 & alpha*delta*power(kappa,2))*U2(m1e,m2e) + \\
 & U1(m1e,m2e)*(alpha*power(kappa,2) - alpha*delta*power(kappa,2) + \\
 & power(alpha,2)*power(kappa,2)*U2(m1e,m2e))))*V1(m1e + 1,m2e + \\
 & 1)))/((1 + kappa*Q1(m1e,m2e))*(1 + kappa*Q2(m1e,m2e))*(1 + \\
 & alpha*U1(m1e,m2e))*(1 + alpha*U2(m1e,m2e)));
 \end{aligned}$$

And that was just one of 6 equations

## Results

$S$	Var	rows	non-zero	dense(%)	Steps	RT (m:s)
20	2400	2568	31536	0.48	5	0 : 03
50	15000	15408	195816	0.08	5	0 : 19
100	60000	60808	781616	0.02	5	1 : 16
200	240000	241608	3123216	0.01	5	5 : 12

Convergence for  $S = 200$

Iteration	Residual
0	1.56(+4)
1	1.06(+1)
2	1.34
3	2.04(-2)
4	1.74(-5)
5	2.97(-11)

## Extensions

Complementarity problems

Continuous time setting